

- 7.1 Two disadvantages associated with database systems are listed below.
- \* setup of the database systems requires more knowledge, money, skills and time.
  - \* The complexity of database may result in poor performance.

- 1.7
1. A file-processing system is more specific to the problem at hand while a DBMS is more general. A file-processing system used by a university is difficult to use in a hospital setting. While a DBMS once written can be used in different places.
  2. It is difficult to ensure atomicity in a conversational file-processing system while it is a lot easier in a DBMS. often wrapping a set of SQL statements in a "BEGIN/TRANSACTIONS" are often enough in the relational DBMS world.
  3. Protecting against concurrent-access anomalies in a file-processing system is difficult, using a DBMS is much easier to protect against concurrent-access anomalies.
  4. most DBMS have a concept of a user and what access that user has. Enforcing such authorization in a file-processing system is really difficult.

- 1.9
1. security - Since DBMS have the concept of a ROLE (user) it easier for setting access management.
  2. Needs to offer atomicity when needed - If atomicity is not provided, inconsistency will be inevitable.
  3. Needs to offer a simple and efficient way to query data.
  4. Needs to offer durability i.e. once an update or an insert has happened it must be persisted.
  5. A DBMS needs to offer a way for protecting against concurrent-access anomalies.

- 1.12
- \* Earlier-generation database applications used a two-tier architecture, where as a three-tier architecture is used by a modern database application.
  - \* In a two-tier architecture the application resides at the client machine, and invokes database system functionality at the server machine through query language statements. In a three-tier architecture the client machine acts as merely a front end and does not contain any direct database calls; the front end communicates with an application server. The application server, in turn, communicates with a database system to access data.
  - \* Three-tier applications provide better security as well as better performance than two-tier applications.

Notes: Even though the book classifies database applications in two, the reality is that most famous applications use four-tier architecture. If we take most chatting mobile applications they are a three-tier architecture with a local database such as SQLite for caching the data and accessing it when the mobile is not connected to the internet.

(2.5) The result attributes include all attribute values of student followed by all attributes of advisor.  
 The rules in the result are as follows: For each student who has an advisor, the result has a row containing the student's attributes, followed by s.id attribute identical to the student's ID attribute, followed by the i.id attribute containing the ID of the student advisor.  
 students who do not have an advisor will not appear in the result. A student who has more than one advisor will appear a corresponding number of times in the result.

- (2.6) a.  $\Pi_{personname} (\sigma_{city = 'miami'} (employee))$   
 b.  $\Pi_{personname} (\sigma_{salary > 100000} (employee \bowtie works))$   
 c.  $\Pi_{personname} (\sigma_{salary > 100000 \wedge city = 'miami'} (employee \bowtie works))$

- (2.7) a.  $\Pi_{branchname} (\sigma_{branch\ city = 'chicago'} (branch))$   
 b.  $\Pi_{ID} (\sigma_{branch\_name = 'Downtown' \wedge loan\_loan\_name = borrower\_loan\_number} (borrower))$

(2.12) a.

Relation Name	Primary Key
branch	branch-name
customer	ID
loan	loan-number
borrower	{ID, loan-number}
account	account-number
depositor	{ID, account-number}

- b.
- | Relation  | Foreign Key  |
|-----------|--|
| branch    | No foreign key   |
| customer  | No foreign key   |
| loan      | branch-name  |
| borrower  | ID - a foreign key referencing customer relation, loan-number - a foreign key referencing loan relation        |
| account   | branch-name  |
| depositor | ID - a foreign key referencing customer relation, account number - a foreign key referencing account relation. |

- (2.14) a.  $\Pi_{ID, personname} (employee \bowtie employee.ID = works.ID (\sigma_{company\ name = 'Big Bank'} (works)))$   
 b.  $\Pi_{ID, personname, city} (employee \bowtie employee.ID = works.ID (\sigma_{company\ name = 'Big Bank'} (works)))$   
 c.  $\Pi_{ID, personname, street, city} (employee \bowtie employee.ID = works.ID (\sigma_{company\ name = 'Big Bank' \wedge salary > 100000} (works)))$   
 d.  $\Pi_{ID, personname} (employee \bowtie employee.ID = works.ID \bowtie works \bowtie works.company\ name = company.company\ name)$

- (2.16)
1. when a value of an attribute is unknown.
  2. when a value of an attribute does not exist.



- 2.18
- a.  $\Pi ID, name (\sigma dept\_name = "Physical" (instructor))$
  - b.  $\Pi ID, name (instructor \bowtie instructor.dept\_name = department.dept\_name (\sigma building = "Watson" (department)))$
  - c.  $\Pi student.ID, student.name (\sigma dept\_name = "Comp.Sci" (student \bowtie student.ID = takes.ID \text{ takes } \bowtie take.course\_id = course.course\_id \text{ course}))$
  - d.  $\Pi student.ID, student.name (\sigma year = 2018 (student \bowtie student.ID = takes.ID \text{ takes}))$
  - e.  $\Pi ID, name (student) - \Pi student.ID, student.name (\sigma year = 2018 (student \bowtie student.ID = takes.ID \text{ takes}))$

- 3.9
- a. 

```
SELECT e.ID, e.person_name, city
FROM employee AS e, works AS w
WHERE w.company_name = 'First Bank Corporation' AND w.ID = e.ID
```
  - b. 

```
SELECT ID, name, city
FROM employee
WHERE ID IN (
    SELECT ID
    FROM works
    WHERE company_name = 'First Bank Corporation' AND salary > 10000
)
```

another solution:

```
SELECT e.ID, e.person_name, city
FROM employee AS e, works AS w
WHERE w.company_name = 'First Bank Corporation' AND w.ID, e.ID
AND w.salary > 10000
```

- c. 

```
SELECT ID
FROM works
WHERE company_name <> 'First Bank Corporation'
```

another solution:

```
SELECT ID
FROM employee
WHERE ID NOT IN (
    SELECT ID
    FROM works
    WHERE company_name = 'First Bank Corporation'
)
```

- d. 

```
SELECT ID
FROM works
WHERE salary > ALL (
    SELECT salary
    FROM works
    WHERE company_name = 'Small Bank Corporation'
)
```

3.10 e. SELECT S.company-name  
 FROM company AS S  
 WHERE NOT EXISTS (  
   (SELECT city  
   FROM company  
   WHERE company-name = 'Small Bank Corporation')  
 )  
 EXCEPT  
 (SELECT city  
   FROM company AS T  
   WHERE T.company-name = S.company-name  
 )  
 )

f. SELECT company-name  
 FROM works  
 GROUP BY company-name  
 HAVING COUNT(DISTINCT ID) >= ALL (  
   SELECT COUNT(DISTINCT ID)  
   FROM works  
   GROUP BY company-name  
 )

g. SELECT company-name  
 FROM works  
 GROUP BY company-name  
 HAVING AVG(salary) > (  
   SELECT AVG(salary)  
   FROM works  
   WHERE company-name = 'First Bank Corporation'  
 )

3.12

- a. INSERT INTO course (course-id, title, dept-name, credits)  
 VALUES ('CS-001', 'Weekly Seminar', 'Comp Sci.', 0);
- b. INSERT INTO section (course-id, sec-id, semester, year) VALUES  
 ('CS-001', '1', 'Fall', 2017)
- c. INSERT INTO takes (id, course-id, sec-id, semester, year)  
 SELECT student-id, 'CS-001', '1', 'Fall', 2017  
 FROM student  
 WHERE student-dept-name = 'Comp. Sci.'

(3.12)

e. DELETE FROM course  
WHERE course-id = 'CS-101';

create table section

course-id varchar(8),  
sec-id varchar(8),  
semester varchar(6),

year check (semester in ('Fall', 'Winter', 'Spring', 'Summer')),  
numeric(4,0) check (year > 1701 and year < 2100),

building varchar(15),

room-number varchar(7),

time-slot-id varchar(4),

primary key (course-id, sec-id, semester, year),

foreign key (course-id) references course (course-id)  
on delete cascade

foreign key (building, room-number) references classroom (building, room-number)  
on delete set null  
);

f. DELETE FROM takes

WHERE course-id IN (

SELECT course-id

FROM course

WHERE LOWER(title) LIKE '%advanced%'

)

(3.15) a. WITH all-branches-in-brooklyn (branchname) AS (

SELECT branch-name

FROM branch

WHERE branch-city = 'Brooklyn'

)

SELECT ID, customer.name

FROM customer AS c1

WHERE NOT EXISTS (

(SELECT branch-name FROM all-branches-in-brooklyn)

EXCEPT

(SELECT branch-name

FROM account INNER JOIN depositor

ON account.account-number = depositor.account-number

WHERE depositor.ID = c1.ID

)

)

b. SELECT SUM(amount)

FROM loan

3.17) SELECT branch-name

FROM branch

WHERE assets > some (

SELECT assets

FROM branch

WHERE branch-city = 'Brooklyn'

);

3.18)

a. SELECT employee.id, employee.person-name

FROM employee INNER JOIN works ON employee.id = works.id

INNER JOIN company ON works.company-name = company-company-name

WHERE employee.city = company.city

b. SELECT E.id, E.person-name

FROM employee AS E INNER JOIN messages ON E.id = manages.id

INNER JOIN employee AS manager\_of\_E ON manages.manager\_id = manager\_of\_E.id

WHERE E.street = manager\_of\_E.street AND

E.city = manager\_of\_E.city;

c. WITH average-salary-per-company (company-name, avg-salary) AS (  
SELECT company-name, AVG(salary)

FROM works

GROUP BY company-name

)  
SELECT E.id, E.person-name

FROM employee AS E INNER JOIN works ON E.id = works.id

WHERE works.salary > (  
SELECT avg-salary

FROM average-salary-per-company

WHERE company-name = works.company-name

)

d. SELECT company-name, SUM(salary) AS total-payroll  
FROM works

GROUP BY company-name

ORDER BY total-payroll ASC

LIMIT 1

3.22)

WHERE UNIQUE (SELECT title FROM course)

WHERE 1 >= ALL (

SELECT COUNT(\*)

FROM course

GROUP BY title

)

WHERE NOT EXISTS (

SELECT \*

FROM course AS c1, course AS c2

WHERE c1.course.id != c2.course.id AND

c1.title = c2.title

)



3.73 SELECT dept-name  
 FROM (SELECT dept-name, sum(salary) AS value FROM instructor GROUP BY dept-name) AS dept-total,  
 (SELECT avg(value) AS value FROM (SELECT dept-name, sum(salary) AS value FROM instructor  
 GROUP BY dept-name) AS x) AS dept-total-avg  
 WHERE dept-total.value >= dept-total-avg.value

3.25 SELECT dept-name  
 FROM department  
 WHERE budget > (SELECT budget FROM department WHERE dept-name = 'philosophy')  
 ORDER BY dept-name ASC;

3.28 SELECT i.id, name  
 FROM instructor AS i  
 WHERE NOT EXISTS (  
 (SELECT course-id FROM course WHERE dept-name = i.dept-name)  
 EXCEPT  
 (SELECT course-id FROM teaches WHERE teaches.id = i.id)  
 )  
 ORDER BY name ASC

3.37 SELECT i.id, name  
 FROM instructor AS i  
 WHERE 'A' NOT IN (  
 SELECT takes-grade  
 FROM takes INNER JOIN teaches  
 ON (takes.course-id, takes.sec-id, takes.semester, takes.year) =  
 (teaches.course-id, teaches.sec-id, teaches.semester, teaches.year)  
 WHERE teaches.id = i.id  
 )